

Greedy Dissection Method for Shared Parallelism in Incomplete Factorization within INMOST Platform

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INMOST

INMOST (<u>www.inmost.org</u>, <u>www.inmost.ru</u>) is a short for:

Integrated

Numerical

Modeling and

Object-oriented

Supercomputing

Technologies

- Distributed meshes
- Distributed linear system assembly
- Parallel linear solver
- Automatic differentiation
- Nonlinear system assembly
- Coupling of unknowns and models

First version during 2012 internship at ExxonMobil

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Parallel Finite Volume Computation on General Meshes

Deringer

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Grids





Dynamic grids



OctreeCutcell example in INMOST-Graphics repository



AdaptiveMesh example for general grid adaptation









Parmetis_AdaptiveRepart

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Models





Jacobian Structure in Multiphysics Problem

Model decomposition for reservoir simulator:





Large systems

• Navier-Stokes system:

$$\frac{\partial \rho \boldsymbol{u}}{\partial t} + \operatorname{div}(\rho \boldsymbol{u} \boldsymbol{u}^T - \mu \nabla \boldsymbol{u} + p \mathbb{I}) = -\frac{\mu}{K_f} \boldsymbol{u},$$
$$\operatorname{div}(\rho \boldsymbol{u}) = 0,$$

• Prothrombin (II):

$$\frac{\partial P}{\partial t} + \operatorname{div}(P\boldsymbol{u} - D\nabla P) = -(k_1\phi_c + k_2B_a + k_3T + k_4T^2 + k_5T^3)P,$$

• Thrombin (IIa):

$$\frac{\partial T}{\partial t} + \operatorname{div}(T\boldsymbol{u} - D\nabla T) = (k_1\phi_c + k_2B_a + k_3T + k_4T^2 + k_5T^3)P - k_6AT,$$

• Clot formation factors (IXa, Xa):

$$\frac{\partial B_a}{\partial t} + \operatorname{div}(B_a \boldsymbol{u} - D\nabla B_a) = (k_7 \phi_c + k_8 T)(B_0 - B_a) - k_9 A B_a,$$

Antithrombin (ATIII):

$$\frac{\partial A}{\partial t} + \operatorname{div}(A\boldsymbol{u} - D\nabla A) = -k_6 A - k_9 A B_a,$$

• Fibrinogen (I):

$$\frac{\partial F_g}{\partial t} + \operatorname{div}(F_g \boldsymbol{u} - D\nabla F_g) = -\frac{k_{10}TF_g}{K_{10} + F_g}$$

Fibrin (la):

$$\frac{\partial F}{\partial t} + \operatorname{div}(F\boldsymbol{u} - D\nabla F) = \frac{k_{10}TF_g}{K_{10} + F_g} - k_{11}F_g$$



Fibrin-polymer:

$$\frac{\partial F_p}{\partial t} = k_{11}F,$$

Inactivated platelets:

$$\frac{\partial \phi_f}{\partial t} + \operatorname{div} \left(k \left(\phi_c, \phi_f \right) \left(\phi_f \boldsymbol{u} - D_p \nabla \phi_f \right) \right) = (k_{12} T - k_{13} \phi_c) \phi_f,$$

Activated platelets:

$$\frac{\partial \phi_c}{\partial t} + \operatorname{div} \left(k \big(\phi_c, \phi_f \big) \big(\phi_c \boldsymbol{u} - D_p \nabla \phi_c \big) \right) = -(k_{12}T - k_{13}\phi_c) \phi_f,$$

Platelets mobility:

$$k(\phi_c, \phi_f) = \tanh\left(\pi\left(1 - \frac{\phi_c + \phi_f}{\phi_{max}}\right)\right)$$

Permeability :

$$\frac{1}{K_f} = \frac{16}{\alpha^2} \phi_p^{\frac{3}{2}} (1 + 56\phi_p) \frac{\phi_{max} + \phi_c}{\phi_{max} - \phi_c}, \qquad \phi_p = \min\left(\frac{7}{10}, \frac{F_p}{7000}\right).$$

Bouchnita, A., Terekhov, K., Nony, P., Vassilevski, Y., & Volpert, V.: A mathematical model to quantify the effects of platelet count, shear rate, and injury size on the initiation of blood coagulation under venous flow conditions. PloS one, 15(7), e0235392, 2020

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Large problems

- Suitable for large problem solutions:
 - Black oil problem (previous talk)
 - 3x unknowns per cell
 - 100M and 200M cells (320 cores, INM RAS cluster):

Ca	se	T_{mat}	T_{prec}	T_{iter}	T_{sol}	T_{upd}	N_n	N_l	Y
SPE10									
SPE10	200M	29.6	34.7	64.1	107.5	0.38	428	3.96	

• Memory per core (significant issue!):

Ca				M_{prec}	
SPE10_	_100M	856.7	165.6	558.4	1943.6
SPE10	_200M	1624	346.5	1054.3	4365.6

 Scaled up to **1B of cells** on 9600 Cray cores by Ahmad Abushaika and Longlong Li at HBKU, Qatar.





Large problems

- Suitable for large problem solutions:
 - Poroelasticity problem
 - 4x unknowns per cell
 - **1.2M** cells (INM RAS cluster, Lomonosov):

Machine	N_{proc}	T_{tot}	T_{asm}	T_{prec}	T_{iter}	T_{upd}
	100	15079.4	1119.8	7245.2	4463	479.7
INM RAS cluster	200	8791.2	582.9	3926.2	2800.9	252.4
	400	4637	300.3	1965.6	1374.2	127
Lomonosov supercomputer	700	3536	234.1	1071.1	1112.42	70.5

• **Target**: full-field geomechanical well stability test.







Flagship INMOST Linear Solver

- **Preconditioned BiCGStab(I)** method¹.
- Preconditioner MPI-parallelization using Additive Schwarz Method.
- Preconditioner OpenMP-parallelization using Bordered Block-Diagonal Form^{2,3}. (New!)
- Multi-level preconditioner with deferred pivoting based on the second-order Crout-ILU^{4,5}.
- Condition estimation of the inverse factors determines the coarse system and tunes dropping tolerances^{6,7}.
- Scaling and reordering of the local system before factorization^{8,9,10}.



References

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- 3) Duff, I. S., Scott, J. A.: *Stabilized bordered block diagonal forms for parallel sparse solvers*. Parallel Computing, 31.3-4 (2005): 275-289. (Bordered block-diagonal form)
- 4) Li N., Saad Y., Chow E.: Crout versions of ILU for general sparse matrices. SIAM Journal on Scientific Computing 25.2 (2003): 716-728. (Crout-ILU)
- 5) Kaporin, I.E.: *High quality preconditioning of a general symmetric positive definite matrix based on its UTU+ UTR+ RTU-decomposition*. Numerical linear algebra with applications 5.6 (1998): 483-509. (Second-order ILU)
- 6) Bollhöfer, M.: *A robust ILU with pivoting based on monitoring the growth of the inverse factors*. Linear Algebra and its Applications 338.1-3 (2001): 201-218. (Tuning dropping tolerances)
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- 9) Olschowka, M., Arnold N.: *A new pivoting strategy for Gaussian elimination*. Linear Algebra and its Applications 240 (1996): 131-151. (Maximizing diagonal product)
- 10) Kaporin, I.E.: *Scaling, reordering, and diagonal pivoting in ILU preconditionings*. Russian Journal of Numerical Analysis and Mathematical Modelling 22.4 (2007): 341-375. (Rescaling for condition reduction)

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Additive Schwarz Method

- For MPI-parallelization.
- Global matrix is composed of local blocks.
- Extend blocks to localize the connection.
- Restricted version.





Distributed system

- Local partition outlier
- Remote partition outlier
- Local partition
- Remote partitions



Bordered Block-Diagonal Form

- For OpenMP-parallelization.
- The problem is to find minimum separator for the K-way partitioning. (NP-hard)
- All algorithms are approximate:
 - Greedy dissection (this work).
 - Usual graph packages: Metis, Patoh, Mondriaan. (todo)
- Work with each block in parallel.
- **Differed pivoting** for the **separator**.





Bordered Block-Diagonal Form

• Greedy algorithm:

- $G graph of A, H = G^{T}G graph product.$
- |G_i| predicts **block size** growth.
- |H_i| predicts **separator** growth.
- Using **priority queue** pick a node with minimal growth in both **block** and **separator**.
- Remove the node from G and H graphs and update **priority queue**.
- Sequential 😕

Algor	ithm 1 Greedy dissection.	
1: fune	ction GreedyDissection $(A, parts)$	
	Let $G = (V, E)$, where	▷ Graph induced by matrix
~	$V = \{1, \dots, n\},$	▷ Vertex set
	$G_i = \{\delta_{ij} a_{ij} \neq 0\}.$	▷ Edge set
	Let $H = G^T G$.	⊳ Graph product
	Let $\boldsymbol{\omega}_B = \{ \boldsymbol{\omega}_B^i = \ G_i\ \forall i \in V \}.$	▷ Weight for block growth
_	Let $\boldsymbol{\omega}_S = \{\omega_H^i = \ S_i\ \forall i \in V\}.$	▷ Weight for separator growth
	Let $B_i = \{\emptyset\}, \forall i \in \{1, \dots, parts\}.$	▷ Weight for separator growth ▷ Blocks
~	Let $S = \emptyset$.	▷ Separator
1.0	Let $R = C = V$.	▷ Candidate row and column sets
	Let $P_i = Q_i = 0, \forall i \in \{1,, n\}.$	Rows and columns reordering matrices
	r = c = q = 1.	Row and column index and current block
	while $R \neq \emptyset$ do	
14:	$i = \arg\min_{j \in R} \left(\omega_B^j + \omega_H^j \right),$	\triangleright Minimal growth of both the block and separator
15:	$R = R \setminus \{i\},$	▷ Remove from the candidate set
16:	$P_i = r,$	p remove nom the candidate bet
17:	r = r + 1,	▷ Enumerate row
18:	$B_q = B_q \cup \{i\},$	> Augment current block
19:	if $i \in S$ then $S = S \setminus \{i\}$.	
20:	for $m \in H_i \cap R$ do	
21:	$\omega_m^S = \omega_m^S - 1.$	> Update weights for separator growth
22:	end for	
23:	for $j \in G_i \cap C$ do	
24:	$C = C \setminus \{j\},\$	
25:	$Q_j = c$,	
26:	c = c + 1.	▷ Enumerate column
27:	for $k \in G_j^T$ do	
28:	$\omega_k^B = \omega_k^B - 1.$	▷ Update weights for block growth
29:	if $k \not\in S$ then	
30:	$S = S \cup \{k\},$	
31:	$\omega_k^S=\omega_k^S-1.$	Add row to separator
32:	$\mathbf{ for } \stackrel{\kappa}{m} \stackrel{\kappa}{\in} H_k \cap R \ \mathbf{ do } \ \omega_m^S = \omega_m^S - 1.$	
33:	$\omega_m^S = \omega_m^S - 1.$	▷ Update weights for separator growth
34:	end for	
35:	end if	
$\frac{36}{37}$:	end for	
38:	end for if $ P > P / nanta then$	
39:	if $ B_q \ge R / parts$ then $R = R \setminus S$,	▷ Remove current separator from candidate set
40:	q = q + 1.	▷ Consider next block
41:	q = q + 1 end if	> Consider next block
	end while	
	for all $i \in S$ do	
44:	$P_i = r,$	
45:	r = r + 1.	
	end for	
_	Q = P.	▷ Match row and column indices.
	return [P,Q,B,S]	
	1 function	













SBBD Schur (no diagonal) 157 的自己 ia. En 18





Ordering and Scaling

- Each successive level is **reordered** and **rescaled**:
 - Reorder to maximize diagonal product. (famous MC64)
 - partially sequential 😕
 - may apply **in parallel** to independent blocks, but **reduce** Schur quality.
 - Rescale into I-dominant matrix.
 - Reorder symmetrically with bandwidth or fill-in reduction algorithm in each block.



Preprocessing methods



Bandwidth reduction

Bandwidth reduction

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Complete blood coagulation problem.

• 13 equations per cell.

- Navier-Stokes problem only.
- 4 equations per cell.
- Fill-in reduction (METIS) is 4 times slower than bandwidth reduction (RCM) algorithm.
- Quite dense matrix graph, big separator.

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Second-order Crout Incomplete LU

- Dual-threshold dropping:
 - **T²** for factorization.
 - **T** for iterations.
- Running condition estimation:





- $\mathbf{\kappa} = \max(|L^{-1}|, |U^{-1}|)$
- **T/K=**CONSt tuning.
- Limit growth of κ .





Transposed matrix traversal:





Schur Complement

- Part that leads to growth of κ is accumulated in C:
 - system reordering after factorization.
- Next level system is the Schur complement:

- $S = C - E (DU)^{-1} D(LD)^{-1} F.$

- Requires forward and backward substitution with sparse right hand side.
- Fill-in control is critical.
- Parallel 🙂

Partially factorized DBBD matrix:





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Analogy to the Algebraic Multigrid

- **Coarse** system should contain the **largest error** of the **smoother**.
- Condition estimation reveals the error in the smoother and provides the coarse-fine splitting of the system.
- Ideal prolongation $P=(-EB^{-1},I)$ and restriction $R=(-FB^{-1},I)^{T}$.
 - (not satisfied by the present method).
- Schur complement corresponds to the **coarse** system.
- Universal but much more computationally complex.
 - (definitely not linear computational complexity)



100x100x100 Poisson problem

Black-oil problem with 1M cells

							t h noo da	1	1	0	4	0	16
threads	1	1	2	4	8	16	threads	1	1	2	4	0	
T_{tot}	42.7	47	29.2 (1.6x)	20.2(2.3x)	13.4 (3.5 x)	10.8 (4.3x)	T_{tot}	46.3	56.6	44.7 (1.3x)	37.2 (1.5x)	30.5 (1.9x)	29.4 (1.9x)
iters	17	23	27	25	27	29	iters	19	30	30	32	29	35
T_{iter}	6	7.8	4.77 (1.6x)	3.2 (2.4x)		1.64(4.8x)	T_{iter}	11.3	25	14.2 (1.8x)	10.7 (2.3x)	6.3 (4x)	6.1 (4.1x)
levels	1	2	5	6	7	7	levels	1	5	6	8	9	10
pivots	0	40564	63588	85689	109995	142988	$_{ m pivots}$	0	167799	201296	278060	355152	458293
T_{prec}	36.2	38.7	23.9 (1.6x)	16.5(2.3x)	10.8 (3.6x)	8.6 (4.5x)	T_{prec}	33.7	30.2	29.1 (1x)	25.2(1.2x)	22.9 (1.3x)	22 (1.4x)
T_{ord}	0.5 (1.4%)	0.7 (2%)	2.6 (11%)	2.4 (15%)	2.3 (21%)	2.55 (29%)	T_{ord}	5.4 (16%)	5.5(18%)	13.1 (45%)	13.3 (53%)	12.3 (54%)	11.5 (52%)
T_{mpt}	0.3 (0.8%)	0.4(1%)	0.4 (1.6%)	0.4(2.4%)	0.41(3.7%)	0.44(5%)	T_{mpt}	2.9 (9%)	2.8(9%)	2.4 (8%)	2.4 (9 %)	2.3 (10%)	2.4 (11%)
T_{snd}	- (-%)	- (-%)	1.9 (8%)	1.9 (11%)	1.8 (16.6%)	2.04(24%)	T_{snd}	- (-%)	- (-%)	7.9 (27%)	8 (32 %)	7.8 (3 4%)	7.8 (36 %)
T_{rcm}	0.2 (0.6%)	0.4(1%)		0.17(1%)	0.1 (0.9%)	0.06(0.8%)	T_{rcm}	2.5 (7%)	2.7 (9%)	2.8 (10%)	2.9 (11%)	2.3 (10%)	1.2 (5.5%)
	0.9 (2.5%)	1 (2.7%)	0.5(2.1%)	0.33(2%)	0.2 (1.9%)	0.17(1.9%)	T_{sc}	4.4 (13%)	4.6 (15%)	2.6 (9%)	1.7 (7%)	1 (4%)	0.8 (4%)
T_{fact}	34.2 (95%)	12.7 (23%)	6.3 (26.4%)	4 (24%)	2.48 (23%)	1.32 (15%)	T_{fact}	20.3 (60%)	9.4 (31%)	5.2 (18%)	3 (12%)	3.1 (13%)	2.4 (11%)
T_{schur}	- (-%)	23.7 (71%)	13.9 (58.2%)	9.2 (55%)	5.35 (49%)	4.04 (47%)	T_{schur}	- (-%)	7.1 (23%)	5.2 (18%)	4.8 (19%)	4.3 (19%)	5.2 (24%)

Single-level algorithm



Conclusions and future directions

- Solved memory but not efficiency issues:
 - Sequential algorithm for DBBD form and Schur growth due to increasing deferred pivoting.
- Future directions:
 - Faster algorithm for **DBBD** form.
 - Parallel maximal product transversal.
 - **Block** elimination, **block** pivoting (*reduce deferred pivoting and Schur growth*).
 - Schur computation using ideal restriction and prolongation (more accurate Schur).
- We are working on a very flexible linear solvers framework S³M to address multiphysics problems:
 - Konshin, I., Terekhov, K.: Sparse System Solution Methods for Complex Problems. In International Conference on Parallel Computing Technologies, (2021, September): 53-73. Springer, Cham.

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Thank you for your attention

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- <u>WWW.INMOST.ORG</u>
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